

FORBIDDING HAMILTON CYCLES IN UNIFORM HYPERGRAPHS

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ABSTRACT. For $1 \leq d \leq \ell < k$, we give a new lower bound for the minimum d -degree threshold that guarantees a Hamilton ℓ -cycle in k -uniform hypergraphs. When $k \geq 4$ and $d < \ell = k - 1$, this bound is larger than the conjectured minimum d -degree threshold for perfect matchings and thus disproves a well-known conjecture of Rödl and Ruciński. Our (simple) construction generalizes a construction of Katona and Kierstead and the space barrier for Hamilton cycles.

1. INTRODUCTION

The study of Hamilton cycles is an important topic in graph theory. A classical result of Dirac [4] states that every graph on $n \geq 3$ vertices with minimum degree $n/2$ contains a Hamilton cycle. In recent years, researchers have worked on extending this theorem to hypergraphs – see recent surveys [16, 18, 26].

To define Hamilton cycles in hypergraphs, we need the following definitions. Given $k \geq 2$, a k -uniform hypergraph (in short, k -graph) consists of a vertex set V and an edge set $E \subseteq \binom{V}{k}$, where every edge is a k -element subset of V . Given a k -graph H with a set S of d vertices (where $1 \leq d \leq k - 1$) we define $\deg_H(S)$ to be the number of edges containing S (the subscript H is omitted if it is clear from the context). The *minimum d -degree* $\delta_d(H)$ of H is the minimum of $\deg_H(S)$ over all d -vertex sets S in H . For $1 \leq \ell \leq k - 1$, a k -graph is called an ℓ -cycle if its vertices can be ordered cyclically such that each of its edges consists of k consecutive vertices and every two consecutive edges (in the natural order of the edges) share exactly ℓ vertices. In k -graphs, a $(k - 1)$ -cycle is often called a *tight* cycle. We say that a k -graph contains a *Hamilton ℓ -cycle* if it contains an ℓ -cycle as a spanning subhypergraph. Note that a Hamilton ℓ -cycle of a k -graph on n vertices contains exactly $n/(k - \ell)$ edges, implying that $k - \ell$ divides n .

Let $1 \leq d, \ell \leq k - 1$. For $n \in (k - \ell)\mathbb{N}$, we define $h_d^\ell(k, n)$ to be the smallest integer h such that every n -vertex k -graph H satisfying $\delta_d(H) \geq h$ contains a Hamilton ℓ -cycle. Note that whenever we write $h_d^\ell(k, n)$, we always assume that $1 \leq d \leq k - 1$. Moreover, we often write $h_d(k, n)$ instead of $h_d^{k-1}(k, n)$ for simplicity. Similarly, for $n \in k\mathbb{N}$, we define $m_d(k, n)$ to be the smallest integer m such that every n -vertex k -graph H satisfying $\delta_d(H) \geq m$ contains a perfect matching. The problem of determining $m_d(k, n)$ has attracted much attention recently and the asymptotic value of $m_d(k, n)$ is conjectured as follows. Note that the $o(1)$ term refers to a function that tends to 0 as $n \rightarrow \infty$ throughout the paper.

Conjecture 1.1. [6, 15] For $1 \leq d \leq k - 1$ and $k \mid n$,

$$m_d(k, n) = \left(\max \left\{ \frac{1}{2}, 1 - \left(1 - \frac{1}{k} \right)^{k-d} \right\} + o(1) \right) \binom{n-d}{k-d}.$$

Conjecture 1.1 has been confirmed [1, 17] for $\min\{k - 4, k/2\} \leq d \leq k - 1$ (the exact values of $m_d(k, n)$ are also known in some cases, e.g., [23, 25]). On the other hand, $h_d^\ell(k, n)$ has also been extensively studied [2, 3, 5, 7, 8, 9, 10, 11, 12, 13, 14, 19, 20, 22, 24]. In particular, Rödl, Ruciński and Szemerédi [20, 22] showed that $h_{k-1}(k, n) = (1/2 + o(1))n$. The same authors proved in [21] that $m_{k-1}(k, n) = (1/2 + o(1))n$ (later they determined $m_{k-1}(k, n)$ exactly [23]). This suggests that the values of $h_d(k, n)$ and $m_d(k, n)$ are closely related and inspires Rödl and Ruciński to make the following conjecture.

Conjecture 1.2. [18, Conjecture 2.18] Let $k \geq 3$ and $1 \leq d \leq k - 2$. Then

$$h_d(k, n) = m_d(k, n) + o(n^{k-d}).$$

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By using the value of $m_d(k, n)$ from Conjecture 1.1, Kühn and Osthus stated this conjecture explicitly for the case $d = 1$.

Conjecture 1.3. [16, Conjecture 5.3] *Let $k \geq 3$. Then*

$$h_1(k, n) = \left(1 - \left(1 - \frac{1}{k}\right)^{k-1} + o(1)\right) \binom{n-1}{k-1}.$$

In this note we provide new lower bounds for $h_d^\ell(k, n)$ when $d \leq \ell$.

Theorem 1.4. *Let $1 \leq d \leq k-1$ and $t = k-d$, then*

$$h_d(k, n) \geq \left(1 - \binom{t}{\lfloor t/2 \rfloor} \frac{[\lfloor t/2 \rfloor]^{[\lfloor t/2 \rfloor]} (\lfloor t/2 \rfloor + 1)^{[\lfloor t/2 \rfloor]}}{(t+1)^t} + o(1)\right) \binom{n}{t}.$$

Theorem 1.5. *Let $1 \leq d \leq \ell \leq k-1$ and $t = k-d$. Then*

$$h_d^\ell(k, n) \geq (1 - b_{t, k-\ell} 2^{-t} + o(1)) \binom{n}{t},$$

where $b_{t, k-\ell}$ equals to the largest sum of $k-\ell$ consecutive binomial coefficients taken from $\{\binom{t}{0}, \dots, \binom{t}{t}\}$.

Theorem 1.4 disproves both Conjectures 1.2 and 1.3.

Corollary 1.6. *For all k ,*

$$h_{k-2}(k, n) \geq \left(\frac{5}{9} + o(1)\right) \binom{n}{2}, \quad h_{k-3}(k, n) \geq \left(\frac{5}{8} + o(1)\right) \binom{n}{3}, \quad h_{k-4}(k, n) \geq \left(\frac{409}{625} + o(1)\right) \binom{n}{4}$$

and in general, for any $1 \leq d \leq k-1$,

$$(1.1) \quad h_d(k, n) > \left(1 - \frac{1}{\sqrt{3(k-d)/2 + 1}}\right) \binom{n}{k-d}.$$

These bounds imply that Conjecture 1.2 is false when $k \geq 4$ and $\min\{k-4, k/2\} \leq d \leq k-2$, and Conjecture 1.3 is false whenever $k \geq 4$.

We will prove Theorem 1.4, Theorem 1.5, and Corollary 1.6 in the next section.

We believe that Conjecture 1.2 is false whenever $k \geq 4$ but due to our limited knowledge on $m_d(k, n)$, we can only disprove Conjecture 1.2 for the cases when $m_d(k, n)$ is known.

This bound $h_{k-2}(k, n) \geq (\frac{5}{9} + o(1)) \binom{n}{2}$ coincides with the value of $m_1(3, n)$ – it was shown in [6] that $m_1(3, n) = (5/9 + o(1)) \binom{n}{2}$, and it was widely believed that $h_1(3, n) = (5/9 + o(1)) \binom{n}{2}$, e.g., see [19]. On the other hand, it is known [17] that $m_2(4, n) = (\frac{1}{2} + o(1)) \binom{n}{2}$, which is smaller than $\frac{5}{9} \binom{n}{2}$. Therefore $k = 4$ and $d = 2$ is the smallest case when Theorem 1.4 disproves Conjecture 1.2. More importantly, (1.1) shows that $h_d(k, n) / \binom{n}{k-d}$ tends to one as $k-d$ tends to ∞ . For example, as k becomes sufficiently large, $h_{k-\ln k}(k, n)$ is close to $\binom{n-d}{k-d}$, the trivial upper bound. In contrast, Conjecture 1.1 suggests that there exists $c > 0$ independent of k and d ($c = 1/e$, where $e = 2.718\dots$, if Conjecture 1.1 is true) such that $m_d(k, n) \leq (1-c) \binom{n-d}{k-d}$.

Similarly, by Theorem 1.5, if $k-\ell = o(\sqrt{t})$, $h_d^\ell(k, n) / \binom{n}{t}$ tends to one as t tends to ∞ because

$$1 - b_{t, k-\ell} 2^{-t} \geq 1 - \frac{k-\ell}{2^t} \binom{t}{\lfloor t/2 \rfloor} \approx 1 - \frac{o(\sqrt{t})}{\sqrt{\pi t/2}}.$$

Theorem 1.5 also implies the following special case: suppose k is odd and $\ell = d = k-2$. Then $t = 2$ and $b_{t, k-\ell} = b_{2, 2} = 3$, and consequently $h_{k-2}^{k-2}(k, n) \geq (\frac{1}{4} + o(1)) \binom{n}{2}$. Previously it was only known that $h_{k-2}^{k-2}(k, n) \geq (1 - (\frac{k}{k+1})^2 + o(1)) \binom{n}{2}$ by (2.1) with $a = \lceil k/(k-\ell) \rceil = (k+1)/2$ from Section 2. When k is large, the bound provided by Theorem 1.5 is much better.

Finally, we do not know if Theorems 1.4 and 1.5 are best possible. Glebov, Person, and Weps [5] gave a general upper bound (far away from our lower bounds)

$$h_d^\ell(k, n) \leq \left(1 - \frac{1}{ck^{3k-3}}\right) \binom{n-d}{k-d},$$

where c is a constant independent of d, ℓ, k, n .

2. THE PROOFS

Before proving our results, it is instructive to recall the so-called *space barrier*. Throughout the paper, we write $X \dot{\cup} Y$ for $X \cup Y$ when sets X, Y are disjoint.

Proposition 2.1. [13] *Let $H = (V, E)$ be an n -vertex k -graph such that $V = X \dot{\cup} Y$ and $E = \{e \in \binom{V}{k} : e \cap X \neq \emptyset\}$. Suppose $|X| < \frac{1}{a(k-\ell)}n$, where $a := \lceil k/(k-\ell) \rceil$, then H does not contain a Hamilton ℓ -cycle.*

A proof of Proposition 2.1 can be found in [13, Proposition 2.2] and is actually included in our proof of Proposition 2.2 below. It is not hard to see that Proposition 2.1 shows that

$$(2.1) \quad h_d^\ell(k, n) \geq \left(1 - \left(1 - \frac{1}{a(k-\ell)}\right)^{k-d} + o(1)\right) \binom{n-d}{k-d}.$$

Now we state our construction for Hamilton cycles – it generalizes the one given by Katona and Kierstead [11, Theorem 3] (where $j = \lfloor k/2 \rfloor$) and the space barrier (where $j = \ell + 1 - k$) simultaneously. The special case with $k = 3, \ell = 2, j = 1$, and $|X| = n/3$ appears in [19, Construction 2].

Proposition 2.2. *Given an integer j such that $\ell + 1 - k \leq j \leq k$, let $H = (V, E)$ be an n -vertex k -graph such that $V = X \dot{\cup} Y$ and $E = \{e \in \binom{V}{k} : |e \cap X| \notin \{j, j+1, \dots, j+k-\ell-1\}\}$. Suppose $\frac{j-1}{a'(k-\ell)}n < |X| < \frac{j+k-\ell}{a(k-\ell)}n$, where $a' := \lfloor k/(k-\ell) \rfloor$ and $a := \lceil k/(k-\ell) \rceil$, then H does not contain a Hamilton ℓ -cycle.*

Proof. Suppose instead, that H contains a Hamilton ℓ -cycle C . Then all edges e of C satisfy $|e \cap X| \notin \{j, j+1, \dots, j+k-\ell-1\}$. We claim that either all edges e of C satisfy $|e \cap X| \leq j-1$ or all edges e of C satisfy $|e \cap X| \geq j+k-\ell$. Otherwise, there must be two consecutive edges e_1, e_2 in C such that $|e_1 \cap X| \leq j-1$ and $|e_2 \cap X| \geq j+k-\ell$. However, since $|e_1 \cap e_2| = \ell$, we have $||e_1 \cap X| - |e_2 \cap X|| \leq k-\ell$, a contradiction.

Observe that every vertex of H is contained in either a or a' edges of C and C contains $\frac{n}{k-\ell}$ edges. This implies that

$$a'|X| \leq \sum_{e \in C} |e \cap X| \leq a|X|.$$

On the other hand, we have $\sum_{e \in C} |e \cap X| < (j-1)\frac{n}{k-\ell}$ or $\sum_{e \in C} |e \cap X| > (j+k-\ell)\frac{n}{k-\ell}$. In either case, we get a contradiction with the assumption $\frac{j-1}{a'(k-\ell)}n < |X| < \frac{j+k-\ell}{a(k-\ell)}n$. \square

Note that by reducing the lower and upper bounds for $|X|$ by small constants, we can conclude that H actually contains no *Hamilton ℓ -path*.

In the proofs of Theorems 1.4 and 1.5, we will consider binomial coefficients $\binom{p}{q}$ with $q < 0$ – in this case $\binom{p}{q} = 0$. We will conveniently write $|X| = xn$, where $0 < x < 1$, instead of $|X| = \lfloor xn \rfloor$ – this does not affect our calculations as n is sufficiently large.

In the proofs of Theorem 1.4 we apply Proposition 2.2 with j and x such that $\frac{j-1}{k} < x < \frac{j+1}{k}$. Note that (when $j \leq k, 1 \leq d \leq k-1$, and $t = k-d$) $\frac{j-1}{k} < x < \frac{j+1}{k}$ implies that $\frac{j-d}{t+1} < x < \frac{j+1}{t+1}$, in particular,

$$(2.2) \quad j-d \leq \lfloor x(t+1) \rfloor \leq j.$$

Proof of Theorem 1.4. Let $x = \lceil t/2 \rceil / (t+1)$.¹ Since $\bigcup_{j=1}^{k-1} (\frac{j-1}{k}, \frac{j+1}{k}) = (0, 1)$ and $1/3 \leq \frac{\lceil t/2 \rceil}{t+1} \leq 1/2$, there exists an integer $j \in [k-1]$ such that $\frac{j-1}{k} < \frac{\lceil t/2 \rceil}{t+1} < \frac{j+1}{k}$. Let $H = (V, E)$ be an n -vertex k -graph such that $V = X \dot{\cup} Y, |X| = xn$ and $E = \{e \in \binom{V}{k} : |e \cap X| \neq j\}$. Since $\frac{j-1}{k}n < |X| < \frac{j+1}{k}n$, H contains no tight Hamilton cycle by Proposition 2.2 with $\ell = k-1$.

Now let us compute $\delta_d(H)$. For $0 \leq i \leq d$, let S_i be any d -vertex subset of V that contains exactly i vertices in X . By the definition of H ,

$$\deg_H(S_i) = \binom{n-d}{t} - \binom{|X|-i}{j-i} \binom{|Y|-(d-i)}{t-j+i}.$$

¹It can be shown that among all $0 < x < 1$ and $0 \leq j \leq k$ such that $\frac{j-1}{k} < x < \frac{j+1}{k}$, our choice of x and j minimizes $\max_{j-d \leq i \leq j} \binom{t}{i} x^i (1-x)^{t-i}$, and thus maximizes $\delta_d(H)$ in (2.3).

Note that this holds for $i > j$ or $i < j - t$ trivially. So we have

$$\begin{aligned}\delta_d(H) &= \min_{0 \leq i \leq d} \left\{ \binom{n-d}{t} - \binom{|X|-i}{j-i} \binom{|Y|-(d-i)}{t-j+i} \right\} \\ &= \binom{n}{t} - \max_{j-d \leq i' \leq j} \left\{ \binom{|X|}{i'} \binom{|Y|}{t-i'} \right\} + o(n^t).\end{aligned}$$

Write $|X| = xn$ and $|Y| = yn$. When $0 \leq i' \leq t$, we have

$$\binom{|X|}{i'} \binom{|Y|}{t-i'} = \frac{(xn)^{i'} (yn)^{t-i'}}{i'!(t-i')!} + o(n^t) = \binom{t}{i'} x^{i'} y^{t-i'} \binom{n}{t} + o(n^t).$$

When $i' < 0$ or $i' > t$, we have $\binom{|X|}{i'} \binom{|Y|}{t-i'} = 0 = \binom{t}{i'} x^{i'} y^{t-i'} \binom{n}{t}$. In all cases, we have

$$(2.3) \quad \delta_d(H) = \binom{n}{t} - \max_{j-d \leq i' \leq j} \left\{ \binom{t}{i'} x^{i'} y^{t-i'} \right\} \binom{n}{t} + o(n^t).$$

Let $a_i := \binom{t}{i} x^i y^{t-i}$. Since $x = \lceil t/2 \rceil / (t+1)$ and $y = 1 - x$, it is easy to see that $\max_{0 \leq i \leq t} a_i = a_{\lceil t/2 \rceil}$ (e.g., by observing $\frac{a_i}{a_{i+1}} = \frac{y}{x} \cdot \frac{i+1}{t-i}$ for $0 \leq i < t$). Moreover, by (2.2), we have $j - d \leq \lceil t/2 \rceil \leq j$. Therefore,

$$\max_{j-d \leq i \leq j} \{a_i\} = a_{\lceil t/2 \rceil} = \binom{t}{\lceil t/2 \rceil} \frac{\lceil t/2 \rceil^{\lceil t/2 \rceil} (\lfloor t/2 \rfloor + 1)^{\lfloor t/2 \rfloor}}{(t+1)^t},$$

and we complete the proof by substituting it into (2.3). \square

Now we turn to the proof of Theorem 1.5, in which we assume that $|X| = n/2$, though a further improvement of the lower bound may be possible by considering other values of $|X|$.

Proof of Theorem 1.5. The proof is similar to the one of Theorem 1.4. Let $H = (V, E)$ be an n -vertex k -graph such that $V = X \dot{\cup} Y$, $|X| = n/2$ and $E = \{e \in \binom{V}{k} : |e \cap X| \notin \{\lceil \ell/2 \rceil, \dots, \lceil \ell/2 \rceil + k - \ell - 1\}\}$. Note that

$$\begin{aligned}a'(k-\ell) &= \left\lfloor \frac{k}{k-\ell} \right\rfloor (k-\ell) \geq k - (k-\ell-1) = \ell + 1 > 2(\lceil \ell/2 \rceil - 1), \text{ and} \\ a(k-\ell) &= \left\lceil \frac{k}{k-\ell} \right\rceil (k-\ell) \leq k + (k-\ell-1) < 2(k - \lfloor \ell/2 \rfloor) = 2(\lceil \ell/2 \rceil + k - \ell).\end{aligned}$$

So we have

$$\frac{\lceil \ell/2 \rceil - 1}{a'(k-\ell)} n < |X| = \frac{n}{2} < \frac{\lceil \ell/2 \rceil + k - \ell}{a(k-\ell)} n.$$

Thus, H contains no Hamilton ℓ -cycle by Proposition 2.2.

Fix $1 \leq d \leq k-1$ and let $t = k-d$. Now we compute $\delta_d(H)$. For $0 \leq i \leq d$, let S_i be any d -vertex subset of V that contains exactly i vertices in X . It is easy to see that

$$\deg_H(S_i) = \binom{n}{t} - \sum_{p=i'}^{i'+k-\ell-1} \binom{|X|}{p} \binom{|Y|}{t-p} + o(n^t),$$

where $i' = \lceil \ell/2 \rceil - i$. Using $|X| = |Y| = n/2$ and the similar calculations in the proof of Theorem 1.4, we get

$$\deg_H(S_i) = \binom{n}{t} - \sum_{p=i'}^{i'+k-\ell-1} \binom{t}{p} \frac{1}{2^t} \binom{n}{t} + o(n^t).$$

By the definition of $b_{t,k-\ell}$, we have

$$\delta_d(H) = \min_{0 \leq i \leq d} \deg_H(S_i) \geq \binom{n}{t} - b_{t,k-\ell} 2^{-t} \binom{n}{t} + o(n^t). \quad \square$$

Corollary 1.6 follows from Theorem 1.4 via simple calculations.

Proof of Corollary 1.6. Let $t = k - d$ and

$$f(t) := \binom{t}{\lfloor t/2 \rfloor} \frac{\lceil t/2 \rceil^{\lceil t/2 \rceil} (\lfloor t/2 \rfloor + 1)^{\lfloor t/2 \rfloor}}{(t+1)^t}.$$

Theorem 1.4 states that $h_{k-t}(k, n) \geq (1 - f(t) + o(1)) \binom{n}{t}$ for any $1 \leq t \leq k - 1$. Since

$$f(2) = \frac{4}{9}, \quad f(3) = \frac{3}{8}, \quad \text{and} \quad f(4) = \frac{216}{625},$$

the bounds for $h_{k-t}(k, n)$, $t = 2, 3, 4$, are immediate. To see (1.1), it suffices to show that for $t \geq 1$,

$$(2.4) \quad 1 - f(t) > 1 - \frac{1}{\sqrt{3t/2 + 1}}.$$

When t is odd, $\frac{\lceil t/2 \rceil^{\lceil t/2 \rceil} (\lfloor t/2 \rfloor + 1)^{\lfloor t/2 \rfloor}}{(t+1)^t} = 1/2^t$; when t is even, $\lceil t/2 \rceil^{\lceil t/2 \rceil} (\lfloor t/2 \rfloor + 1)^{\lfloor t/2 \rfloor} < (\frac{t+1}{2})^t$. Thus, for all t , we have

$$f(t) \leq \binom{t}{\lfloor t/2 \rfloor} \frac{1}{2^t},$$

where a strict inequality holds for all even t . Now we use the fact $\binom{2m}{m} \leq 2^{2m}/\sqrt{3m+1}$, which holds for all integers $m \geq 1$. Thus, for all even t , we have $f(t) \leq 1/\sqrt{3t/2 + 1}$; for all odd t ,

$$f(t) \leq \binom{t}{\lfloor t/2 \rfloor} \frac{1}{2^t} = \frac{1}{2} \binom{t+1}{\lfloor t/2 \rfloor + 1} \frac{1}{2^t} \leq \frac{1}{\sqrt{3(t+1)/2 + 1}} < \frac{1}{\sqrt{3t/2 + 1}}.$$

Hence $f(t) \leq 1/\sqrt{3t/2 + 1}$ for all $t \geq 1$. Moreover, by the computation above, regardless of the parity of t , the strict inequality always holds and thus (2.4) is proved.

We next show that whenever $k \geq 4$ and $2 \leq t \leq k - 1$,

$$1 - f(t) > \max \left\{ \frac{1}{2}, 1 - \left(1 - \frac{1}{k}\right)^t \right\}.$$

This implies that Conjecture 1.3 fails for $k \geq 4$, and Conjecture 1.2 fails for $k \geq 4$ and $\min\{k - 4, k/2\} \leq d \leq k - 2$ (because $m_d(k, n)/\binom{n}{k-d} = \max \left\{ \frac{1}{2}, 1 - \left(1 - \frac{1}{k}\right)^{k-d} \right\} + o(1)$ in this case). It suffices to show that for $k \geq 4$ and $2 \leq t \leq k - 1$,

$$f(t) < 1/2 \quad \text{and} \quad f(t) < \left(1 - \frac{1}{k}\right)^t.$$

The first inequality immediately follows from (2.4) and $1/\sqrt{3t/2 + 1} \leq 1/2$. For the second inequality, note that

$$f(t) < \frac{1}{\sqrt{3t/2 + 1}} < \frac{1}{e} < \left(1 - \frac{1}{k}\right)^{k-1} \leq \left(1 - \frac{1}{k}\right)^t$$

for all $t \geq 5$. For $t = 2, 3$ and all $k \geq 4$, one can verify $f(t) < (3/4)^t \leq \left(1 - \frac{1}{k}\right)^t$ easily. Also, for $t = 4$ and all $k \geq 5$, we have $f(4) < (4/5)^4 \leq \left(1 - \frac{1}{k}\right)^4$. \square

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